

Accelerating an Edge-Based CFD Solver Using Many-Core Co-Processors

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D. Kuzmin (University of Erlangen-Nuremberg)

CFD-Software: FeatFlow2

- Open-source finite element library with demo applications
- ~ 700.000 lines of Fortran 95 code + external libraries
- available online <http://www.featflow.de/en/software/featflow2.html>

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- classical FE-like assembly of coefficient matrices
- edge-based assembly of vectors and operators
- other loops (flux limiter, MVmult, norms, ...)



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- minimally invasive integration of co-processor support
- edge-based assembly of vectors and operators



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Reuse of application code via meta-programming library (C++/Fortran)

2. Milestone: *CUDA port of time consuming parts*

- minimally invasive integration of co-processor support
- edge-based assembly of vectors and operators



Galerkin finite element schemes $\partial_t U + \nabla \cdot \mathbf{F}(U) = 0$

- Weak formulation

$$\int_{\Omega} W \frac{\partial U}{\partial t} - \nabla W \cdot \mathbf{F}(U) \, d\mathbf{x} + \int_{\Gamma} W \mathbf{n} \cdot \mathbf{F}(U) \, ds = 0, \quad \forall W \in \mathcal{W}$$

- Group representation [Fle83]

$$U(\mathbf{x}, t) \approx \sum_j \varphi_j(\mathbf{x}) U_j(t) \quad \mathbf{F}(U) \approx \sum_j \varphi_j(\mathbf{x}) \mathbf{F}(U_j)$$

- Semi-discrete high-order scheme

$$\sum_j m_{ij} \frac{dU_j}{dt} - \sum_j \mathbf{c}_{ji} \cdot \mathbf{F}_j + \sum_j \mathbf{s}_{ij} \cdot \mathbf{F}_j = 0$$

$$m_{ij} = \int_{\Omega} \varphi_i \varphi_j \, d\mathbf{x} \quad \mathbf{s}_{ij} = \int_{\Gamma} \varphi_i \varphi_j \mathbf{n} \, ds \quad \mathbf{c}_{ji} = \int_{\Omega} \nabla \varphi_i \varphi_j \, d\mathbf{x}$$

Galerkin finite element schemes, cont'd

- Galerkin flux decomposition

$$\sum_j \mathbf{c}_{ij} = 0 \quad \Rightarrow \quad - \sum_j \mathbf{c}_{ji} \cdot \mathbf{F}_j = \sum_{j \neq i} \mathbf{c}_{ij} \cdot \mathbf{F}_i - \mathbf{c}_{ji} \cdot \mathbf{F}_j$$

- Semi-discrete high-order scheme [Ku03]

$$\sum_j \left[m_{ij} \frac{dU_j}{dt} + \mathbf{s}_{ij} \cdot \mathbf{F}_j \right] + \sum_{j \neq i} G_{ij} = 0$$

- efficient edge-based assembly of Galerkin fluxes
- precomputation of coefficient matrices (on CPU) and singular transfer to device memory (low storage requirement on GPU)

Algebraic flux correction, *Kuzmin et al.*

- Semi-discrete low-order scheme

$$\underline{m_i \frac{dU_i}{dt}} + \sum_{j \neq i} G_{ij} + \underline{D_{ij}(U_j - U_i)} = 0 \quad m_i = \sum_j m_{ij}$$

mass lumping

artificial dissipation

- Conservative flux decomposition

$$m_i(U_i^H - U_i^L) = \sum_{j \neq i} m_{ij} \left(\frac{dU_i}{dt} - \frac{dU_j}{dt} \right) + D_{ij}(U_i - U_j)$$

antidiffusive fluxes

- **low-order scheme + limited antidiffusion = high-resolution scheme**
- Parallelization of edge-loops is crucial to achieve high overall efficiency

Outline of solution algorithm

- **Initialization:** Transfer edge-data to global device memory

In every **time step**:

- Transfer solution vector into global device memory
- Assemble rhs vector and transfer back to host memory

In every **nonlinear step**:

- Assemble nonlinear parts of operator and residual vector and transfer to host memory
- Combine with constant contributions on host
- Solve nonlinear problem, update solution, and transfer solution into global device memory

overlap transfers
with computation

Preparation of parallel edge-based assembly

- Edge-coloring of the FE sparsity graph

$c \leq 2\Delta - 1$ greedy algorithm

$\Delta \leq c \leq \Delta + 1$ Vizing's algorithm

- Precompute constant coefficient matrices (classical FE-assembly on CPU) and store them into edge-based data structure (AoS/SoA/mixture)



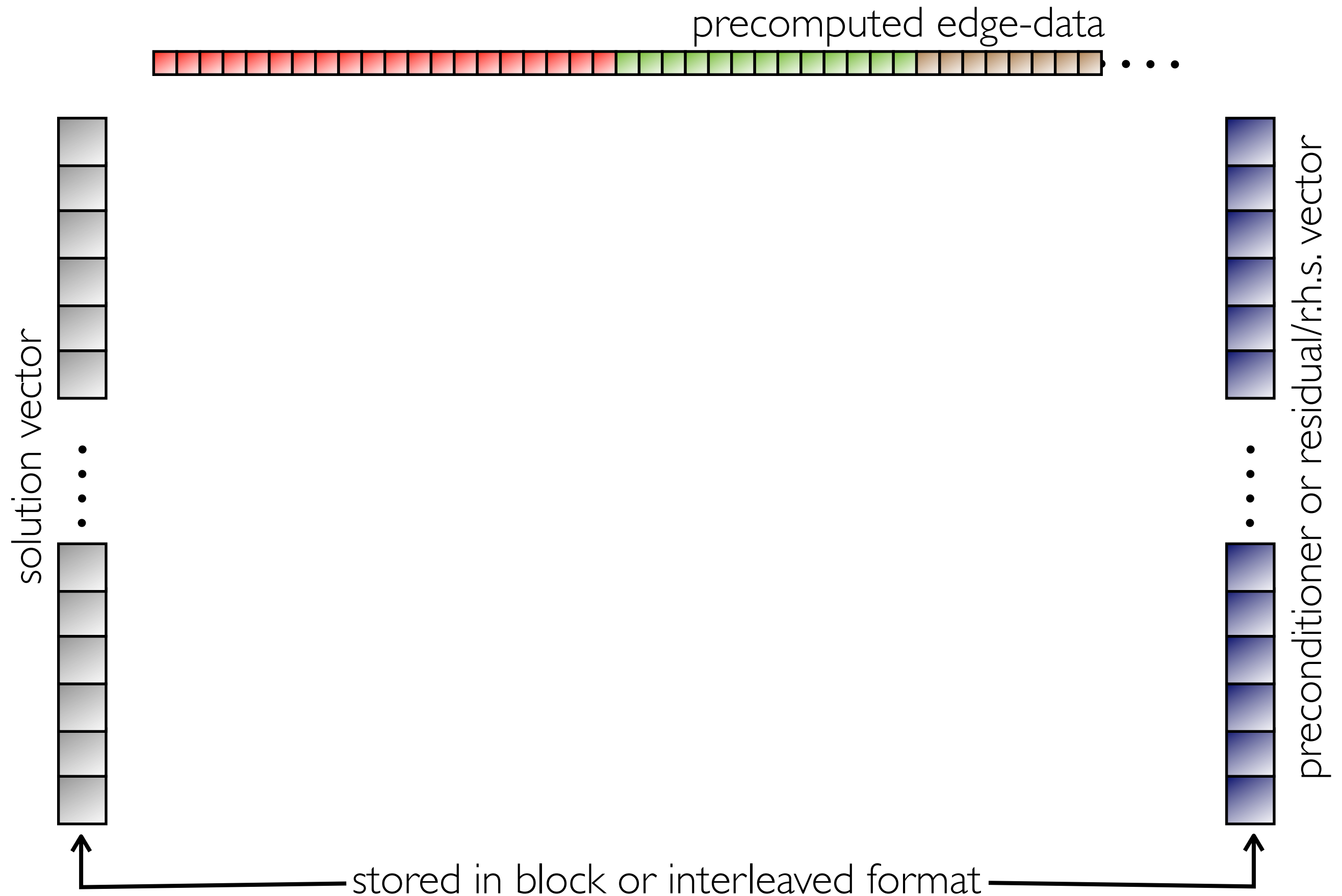
```
struct {  
    int i, j;  
    double m_ij;  
    double cx_ij, cx_ji;  
    double cy_ij, cy_ji;  
} EdgeData[nedge];
```

or

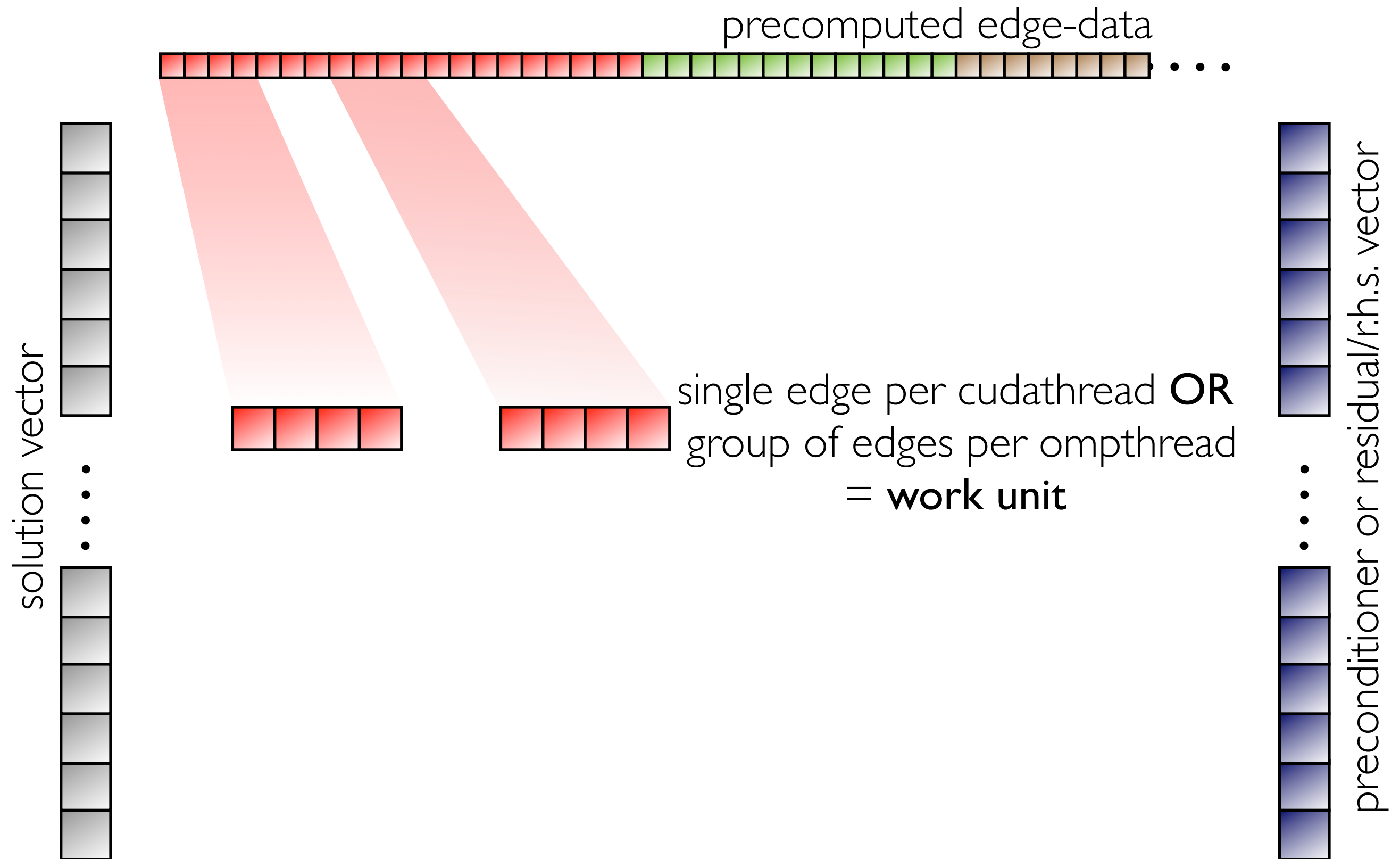
```
int nodes[2][nedge];  
double coeff_m[nedge];  
double coeff_cx[2][nedge];  
double coeff_cy[2][nedge];
```

and index vector separating color groups

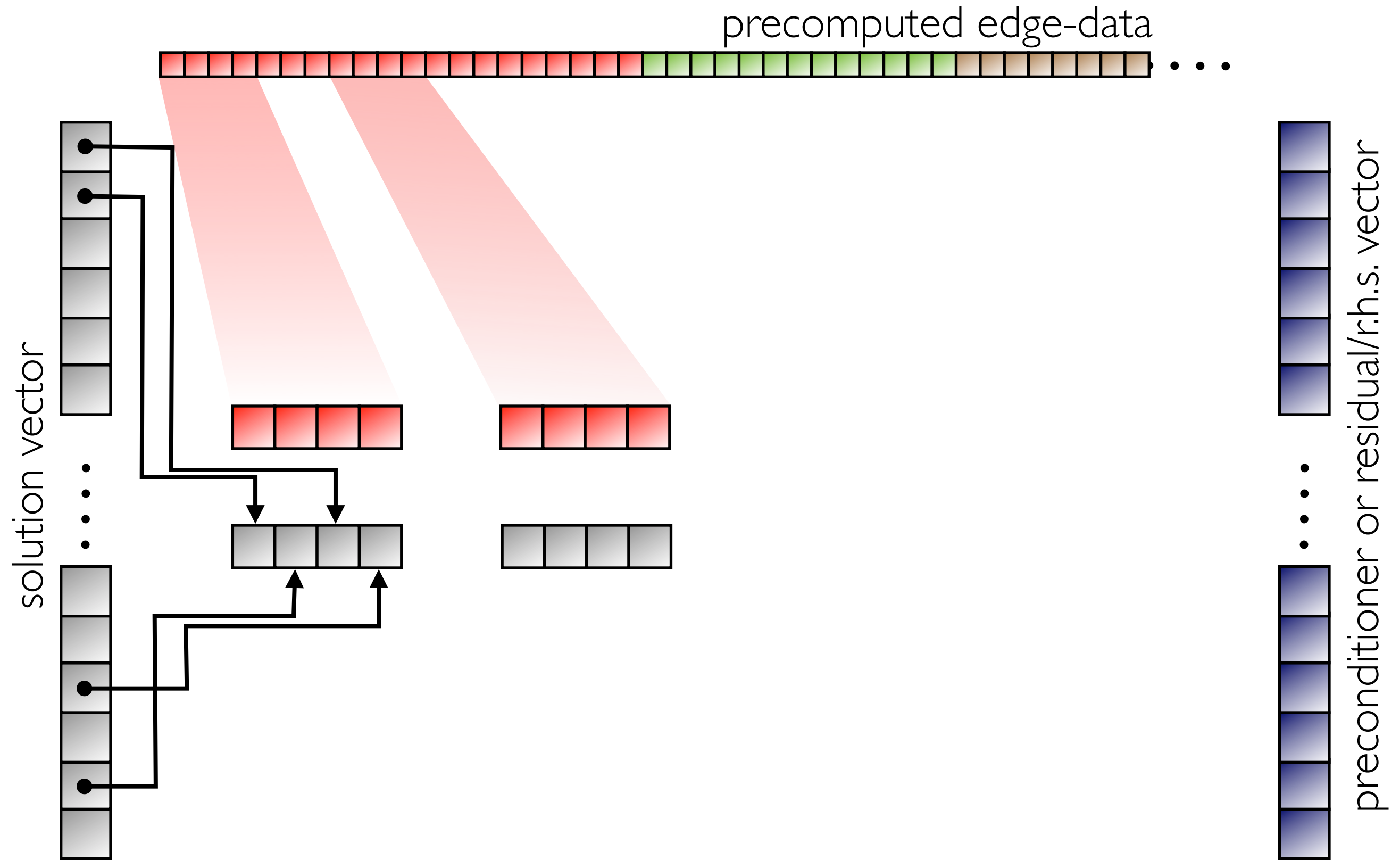
Parallel edge-based assembly



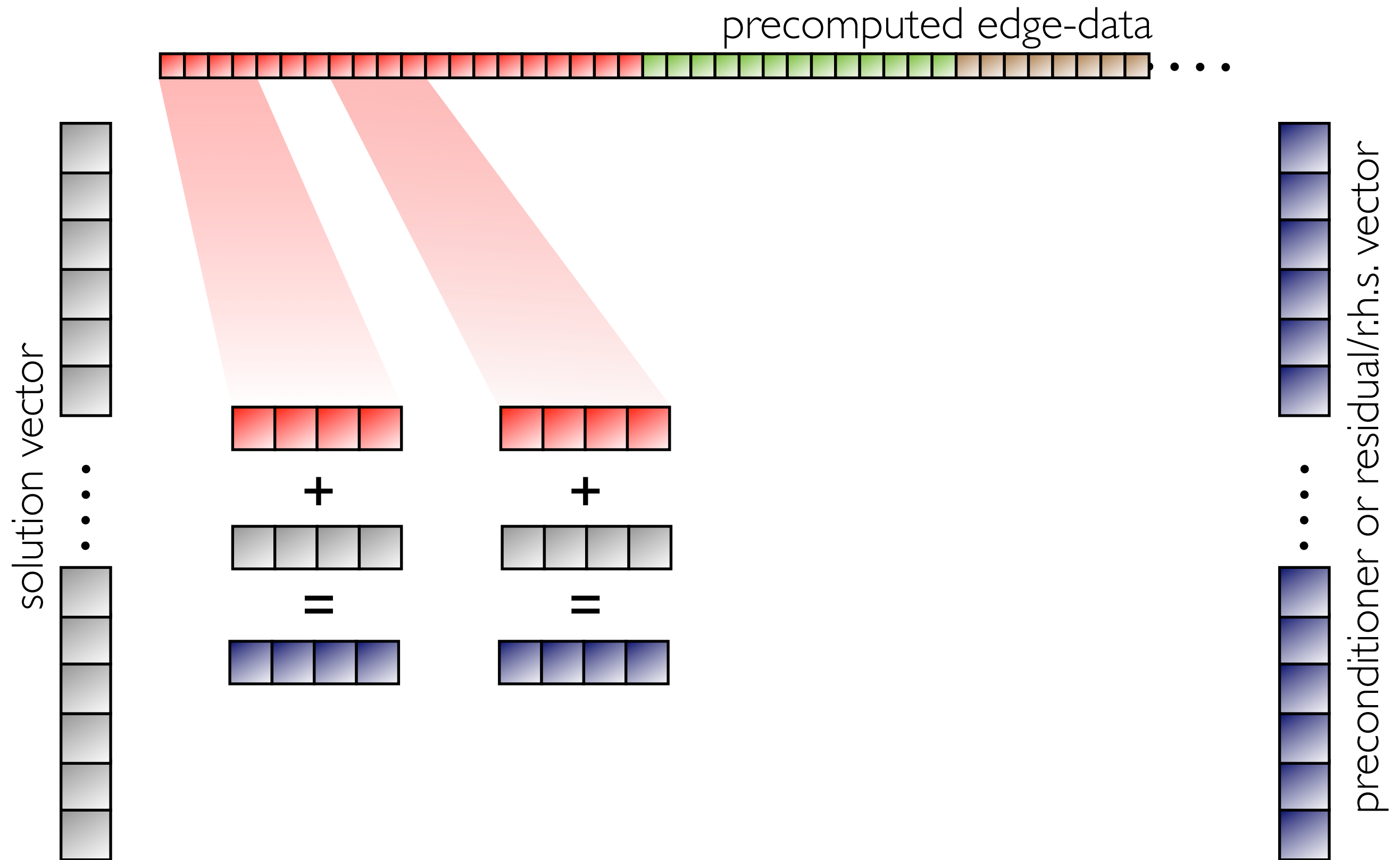
Parallel edge-based assembly



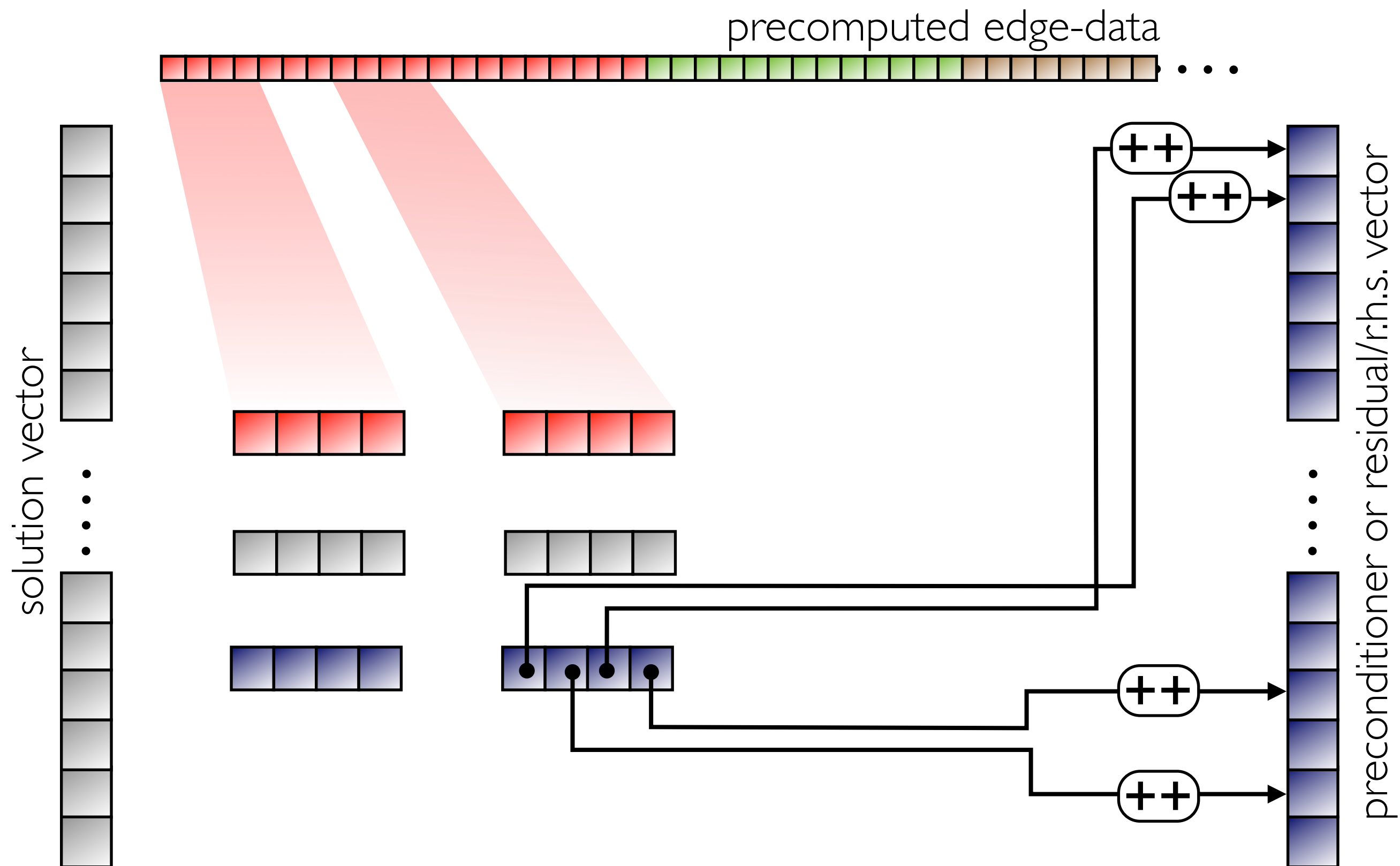
Parallel edge-based assembly



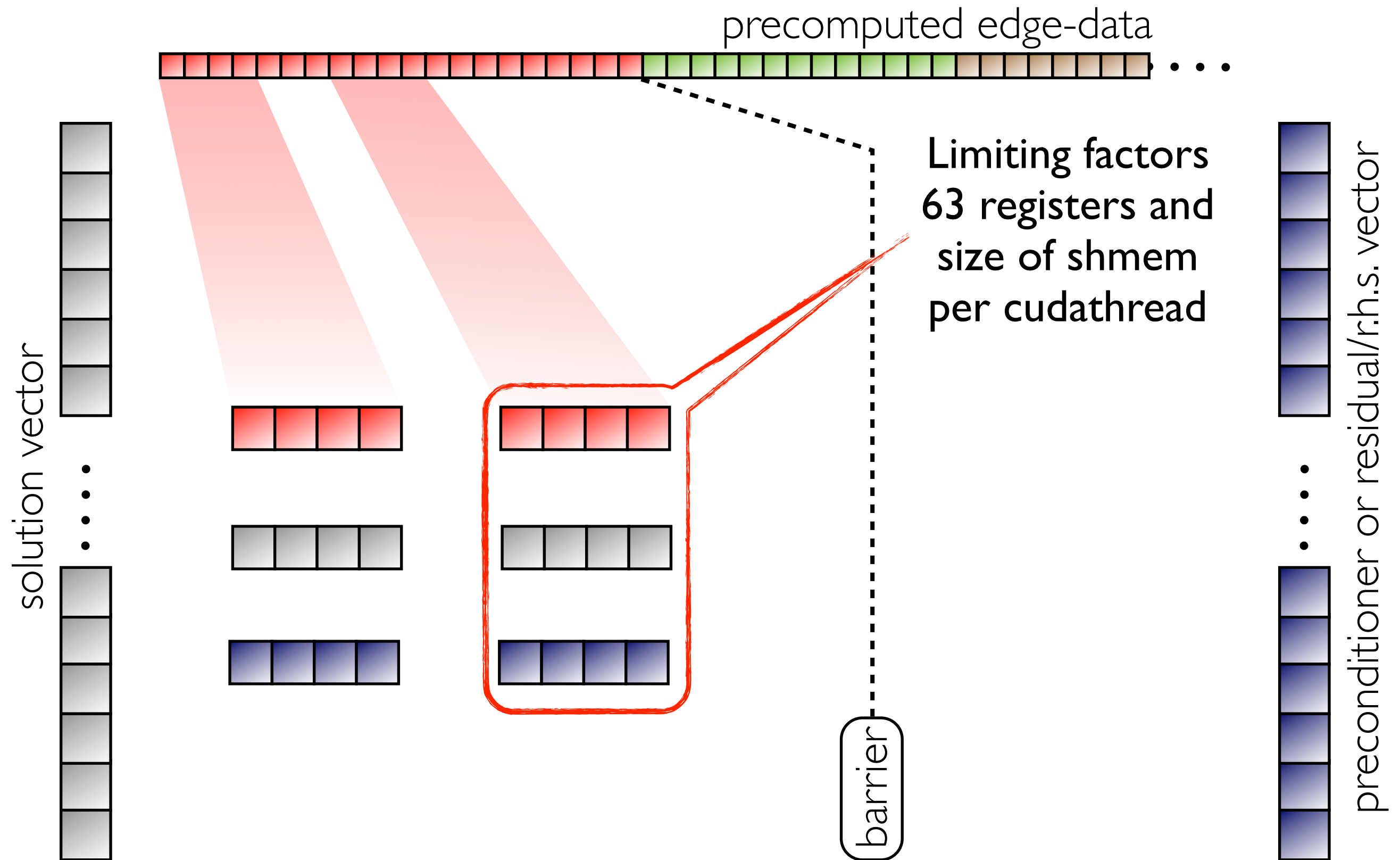
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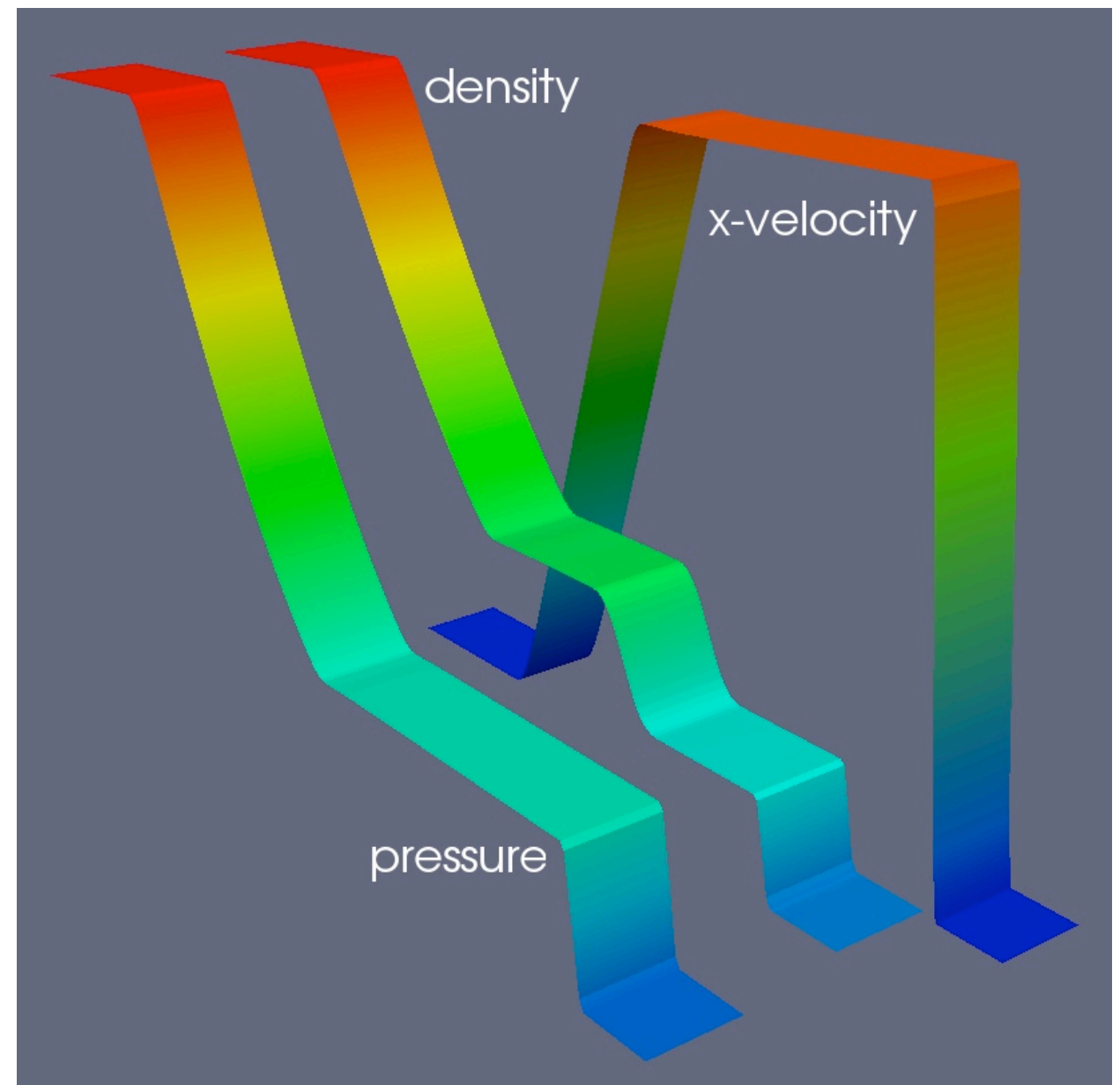


Parallel edge-based assembly



Numerical example

- Sod's Shock tube problem in 2D
- Linearized FEM-FCT (density, pressure)
- Artificial dissipation
 - scalar (39 l.o.c.)
 - Roe-type (55 l.o.c.)
- Q1 finite elements
- Regular grid ($\Delta = 8$)
- Greedy coloring ($c = 14$)
- Gcc 4.4.3, CUDA 4.2






Computational efficiency

Computing platforms

- **P1:** Intel Xeon X5680 at 3.33GHz (2x6, no hyperthreading, 2x12MB L3)
P2: Intel Core i7 at 3.33GHz (1x6, 1x12MB L3) + C2070 (ECC off)
- OpenMP: with 800 edges per „parallel block“
CUDA: with 64 threads per CUDA block

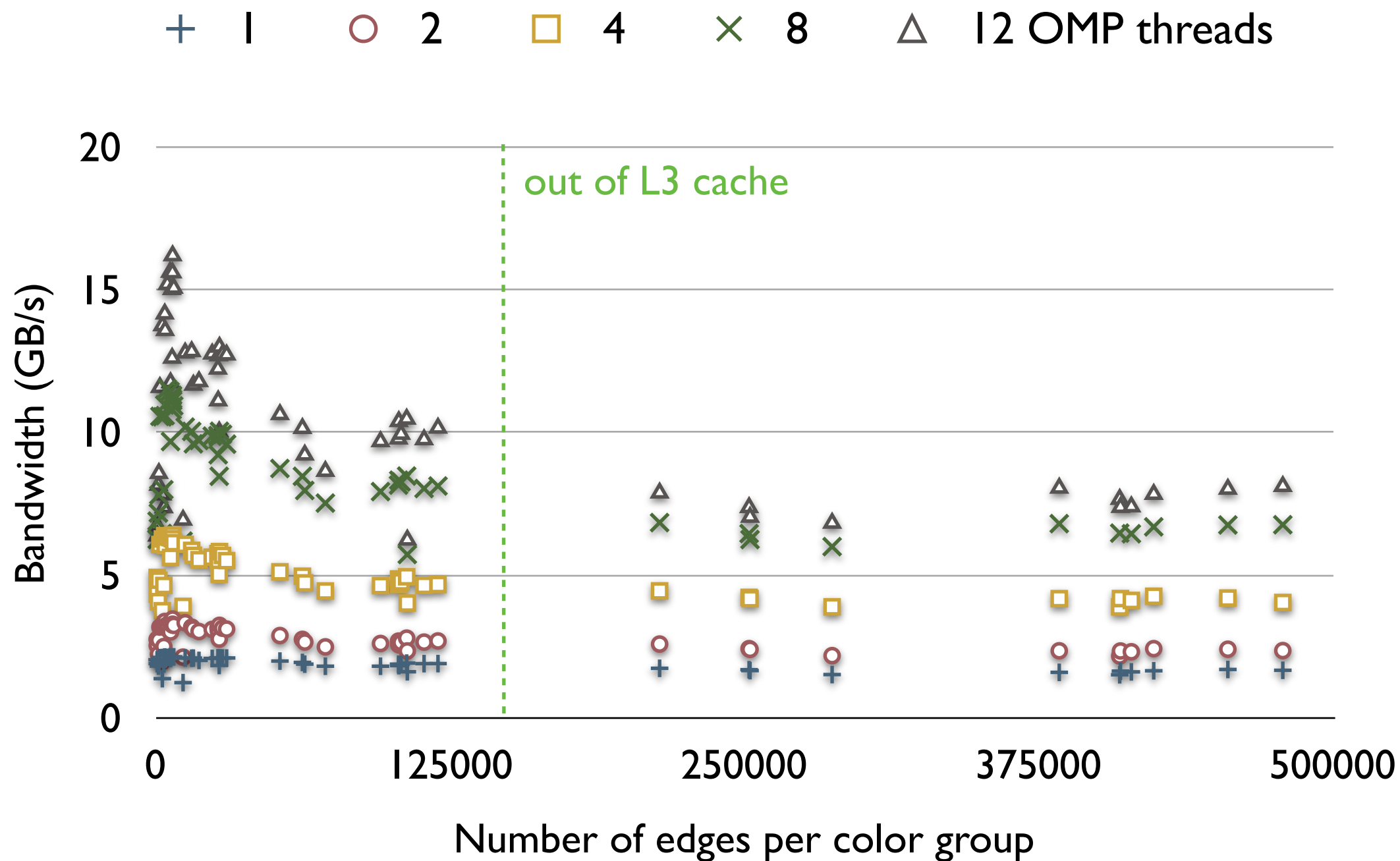
Comparisons

-  micro benchmark: P1-OpenMP vs. P2-CUDA
edge-based vector assembly of a *single* color group
-  meso benchmark: P2-OpenMP vs. P2-CUDA
edge-based vector assembly over all color groups
-  macro benchmark: P2-OpenMP vs. P2-CUDA
„full“ simulation (100 time steps) w/o I/O-operations

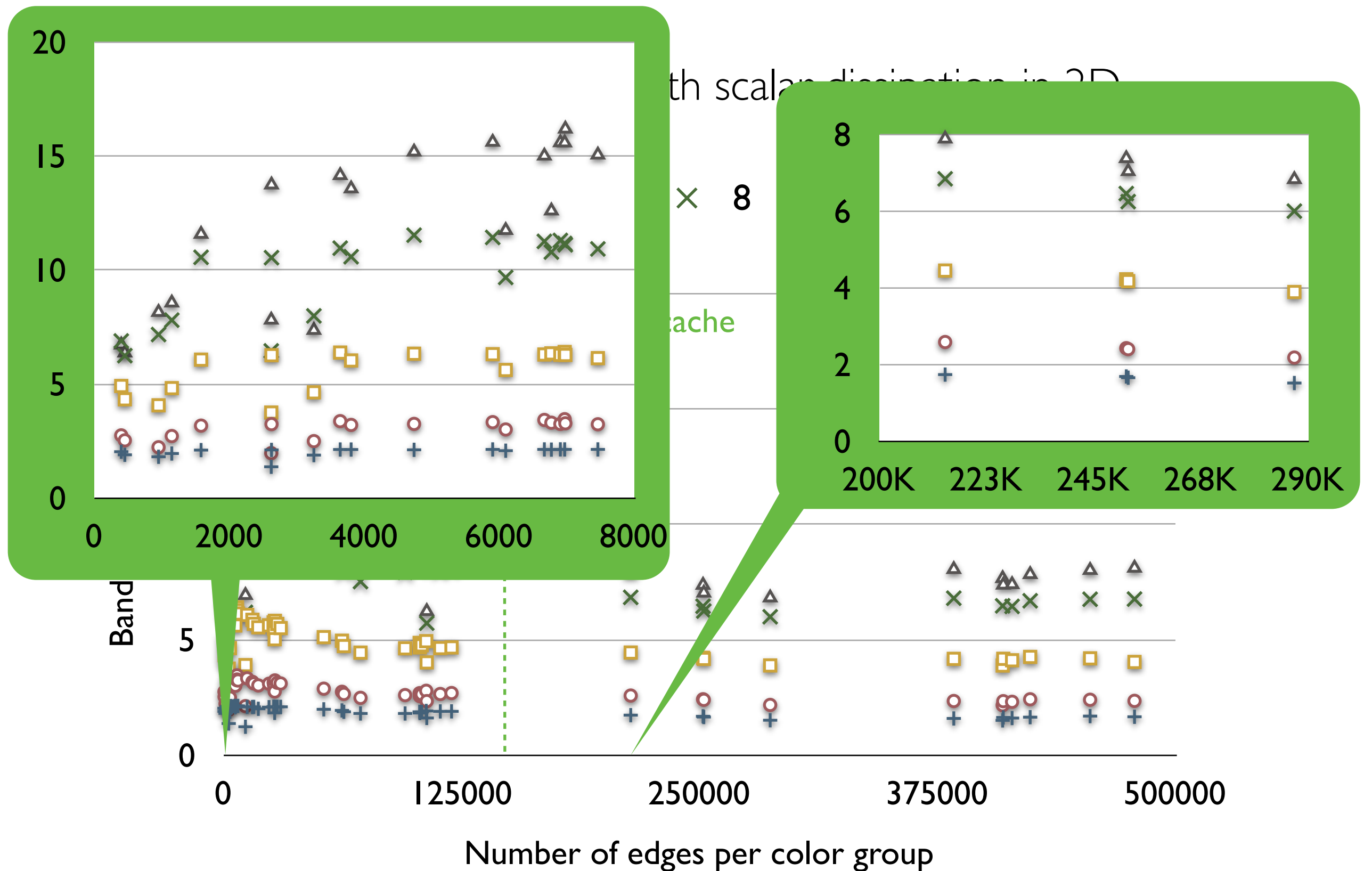
Bandwidth: CPU implementation (PI)



Kernel: inviscid fluxes with scalar dissipation in 2D



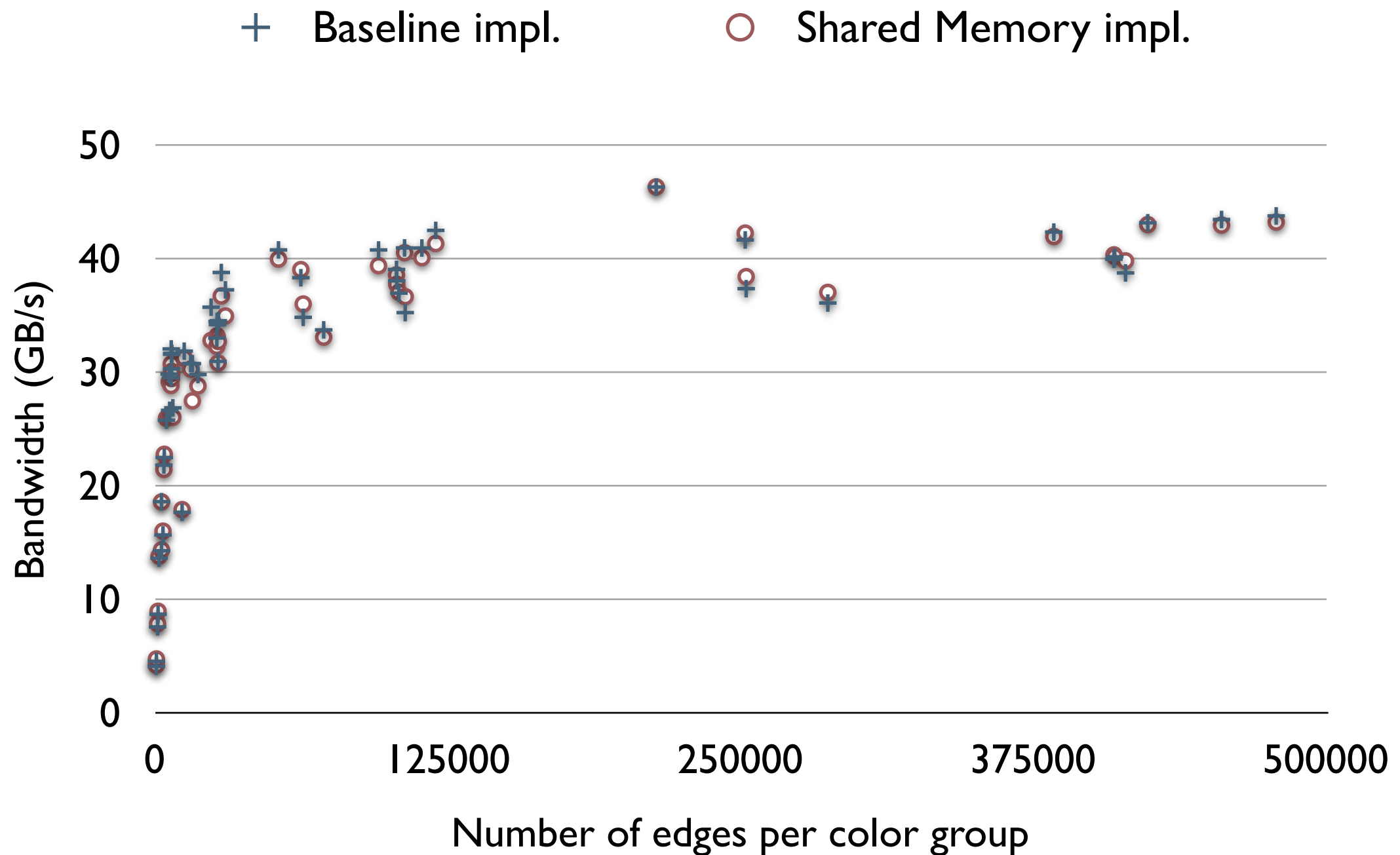
Bandwidth: CPU implementation (PI)



Bandwidth: GPU implementation (P2)



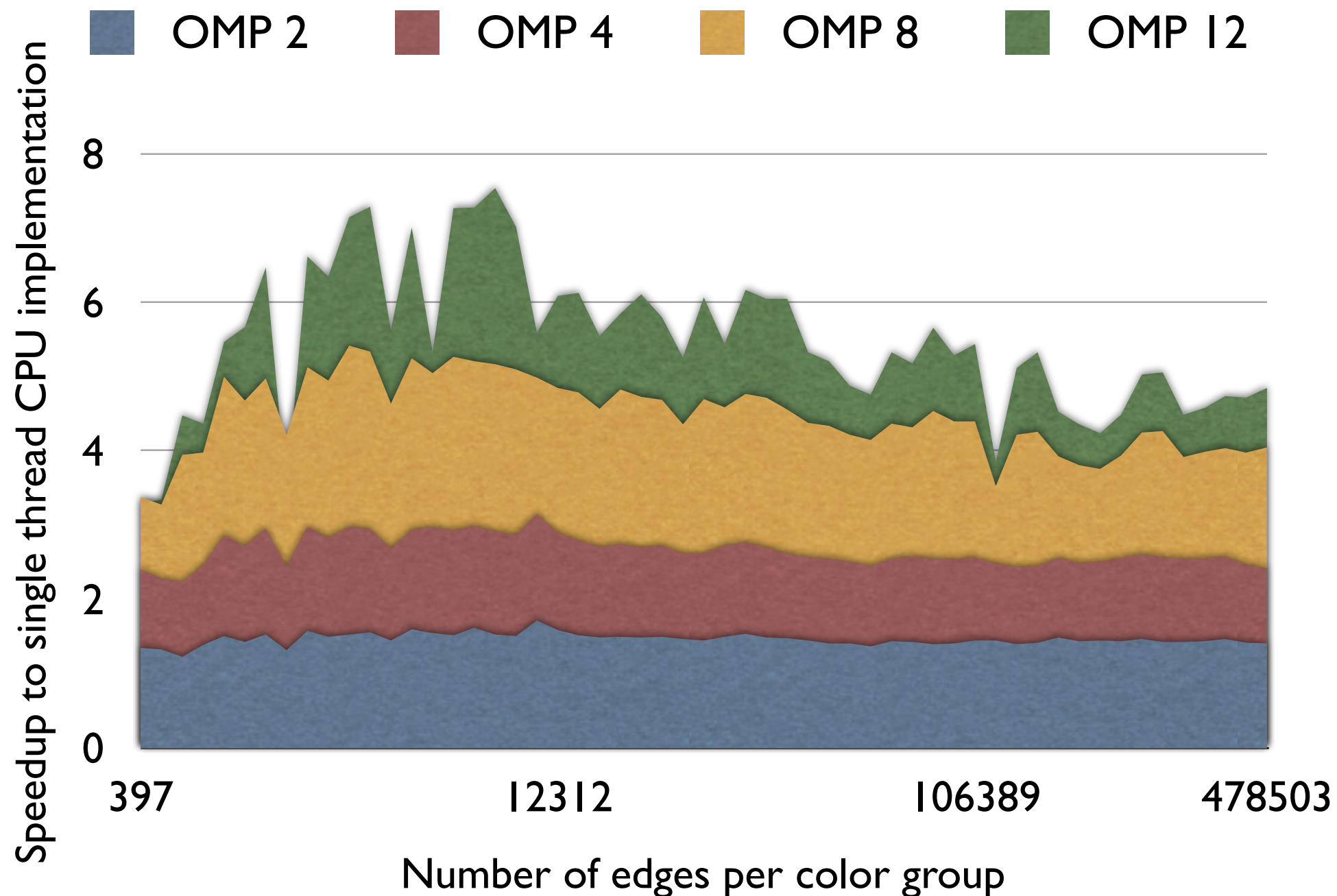
Kernel: inviscid fluxes with scalar dissipation in 2D



Computing time: CPU implementation (PI)



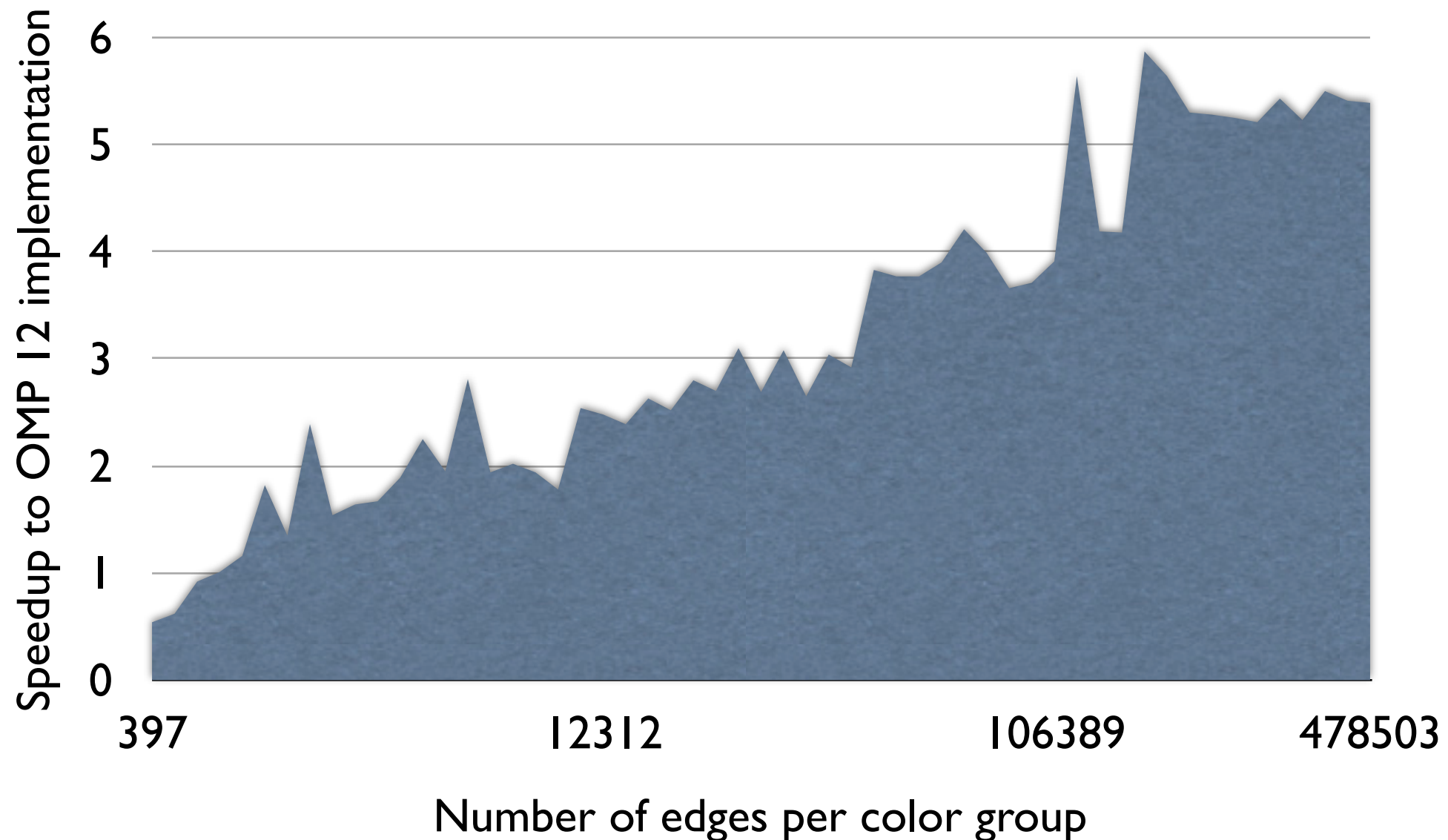
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Computing time: GPU implementation (P2)



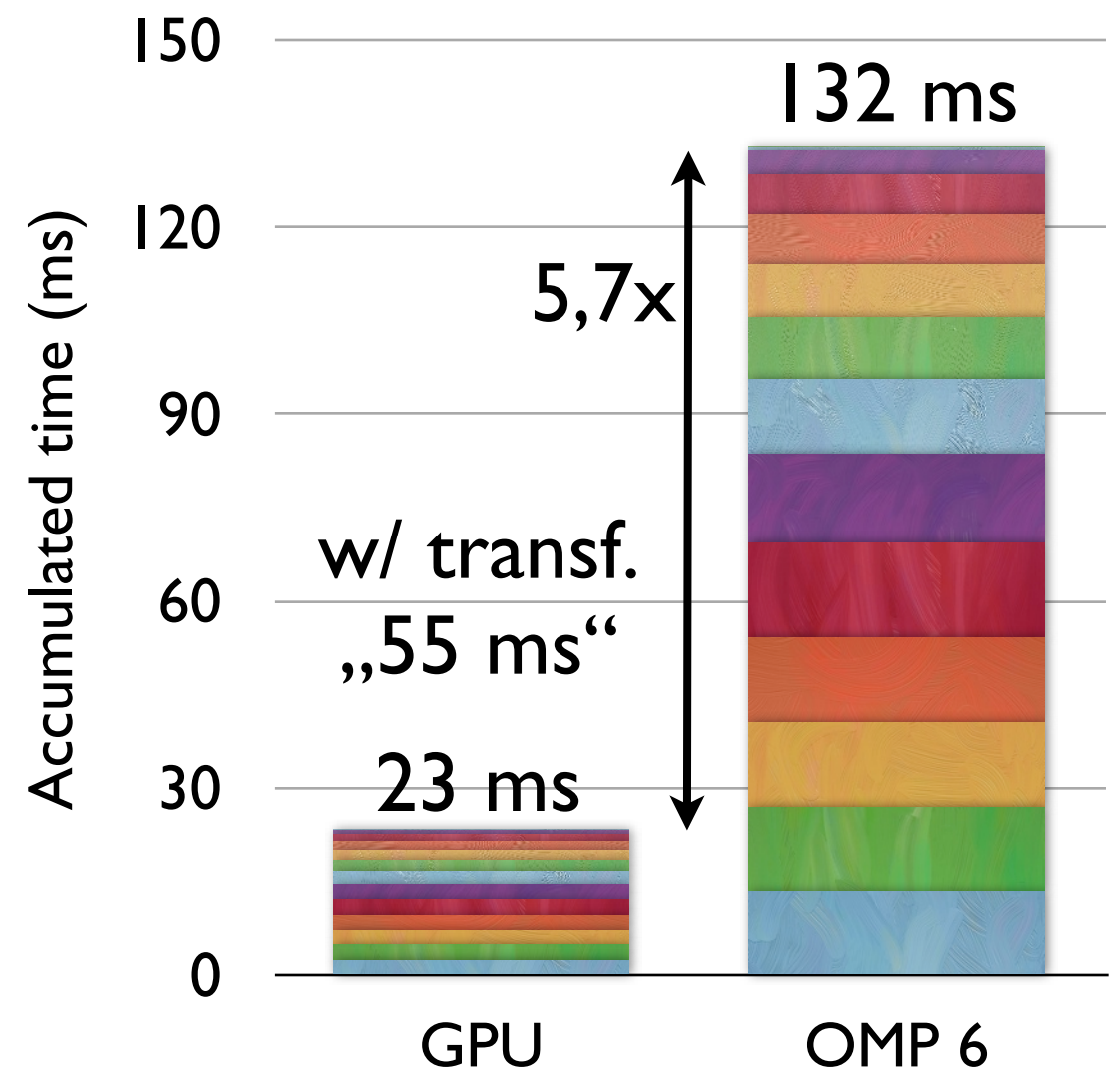
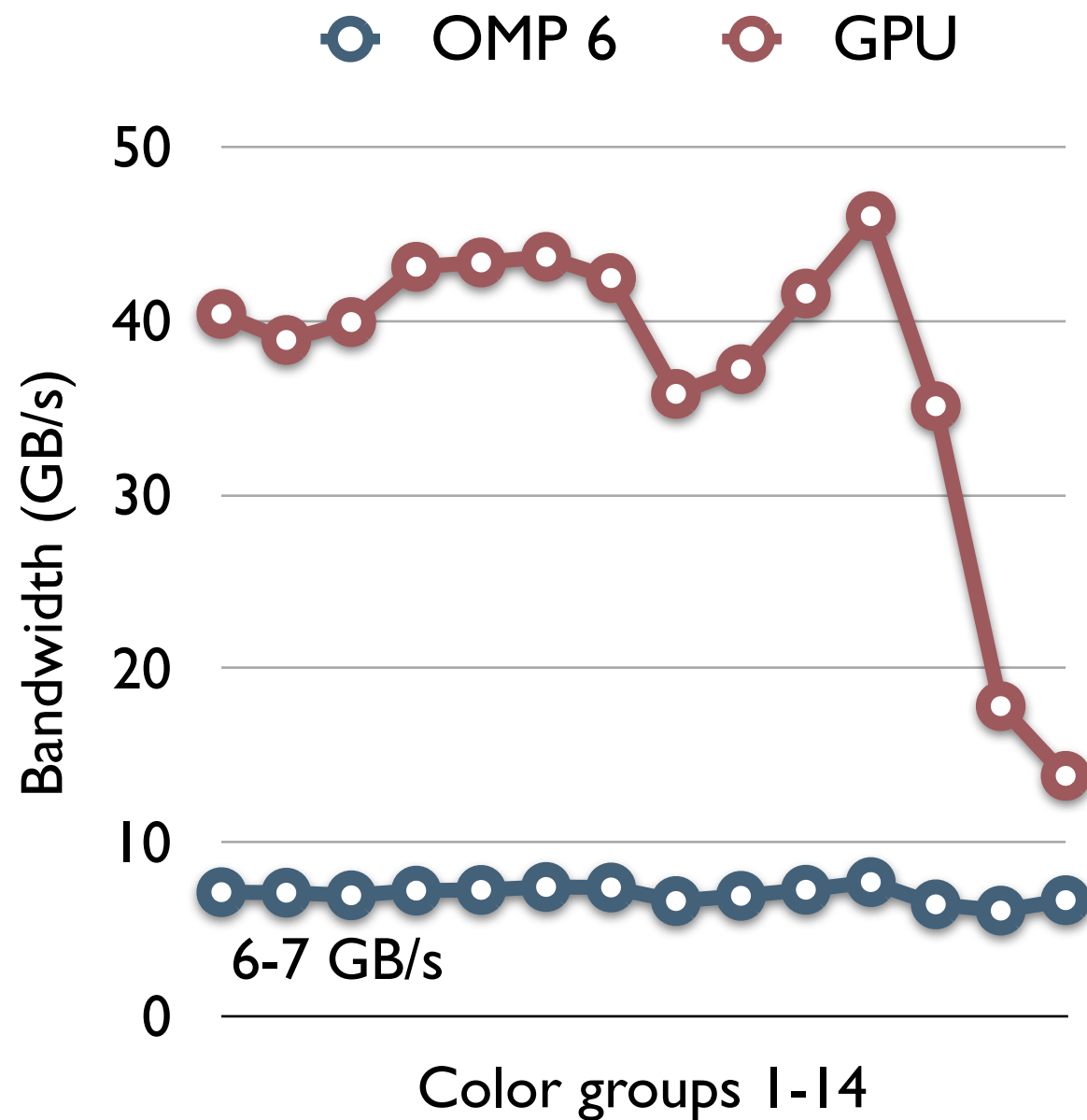
GPU is $>5x$ faster than best CPU implementation



Efficiency of vector assembly (P2)



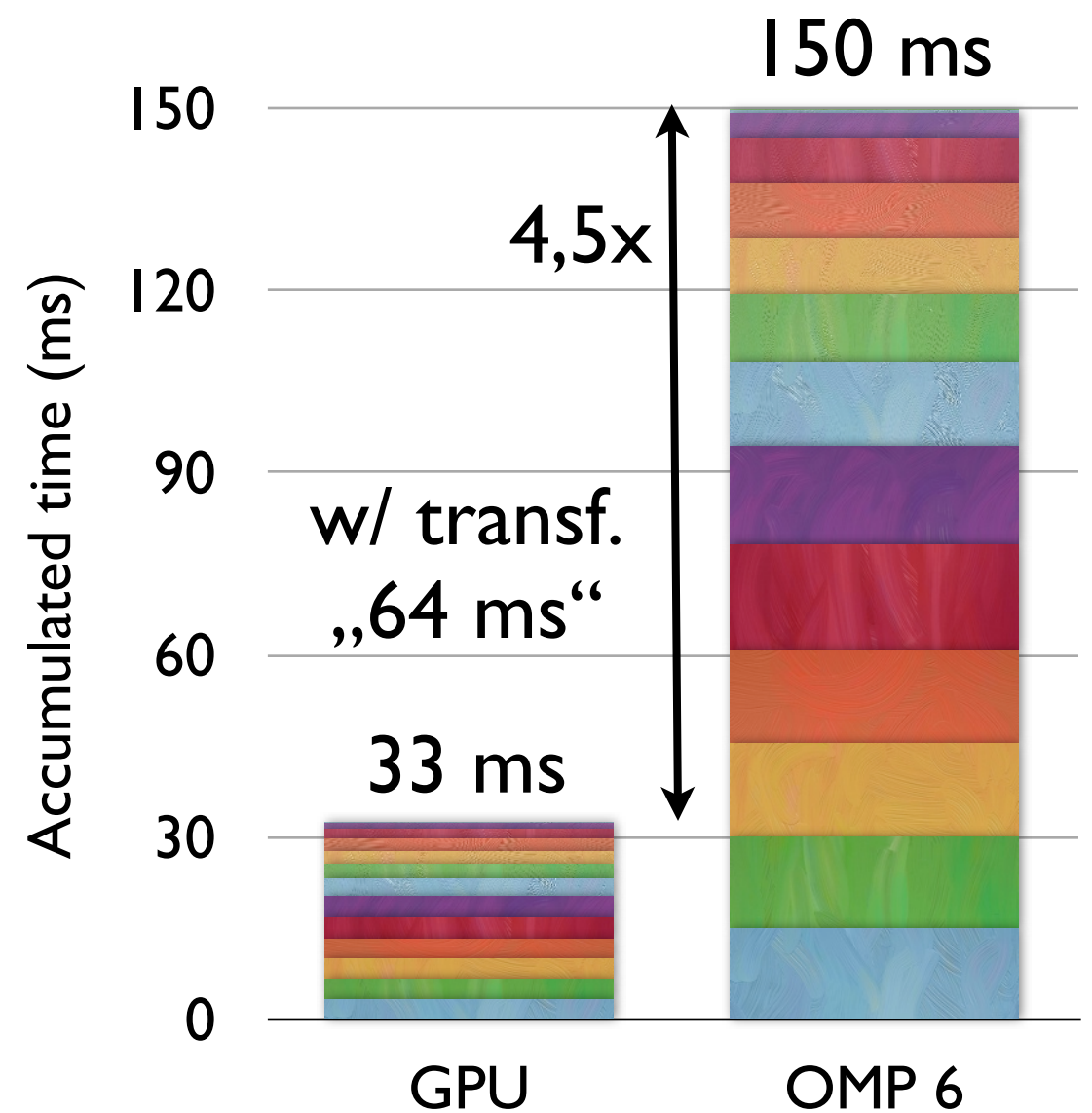
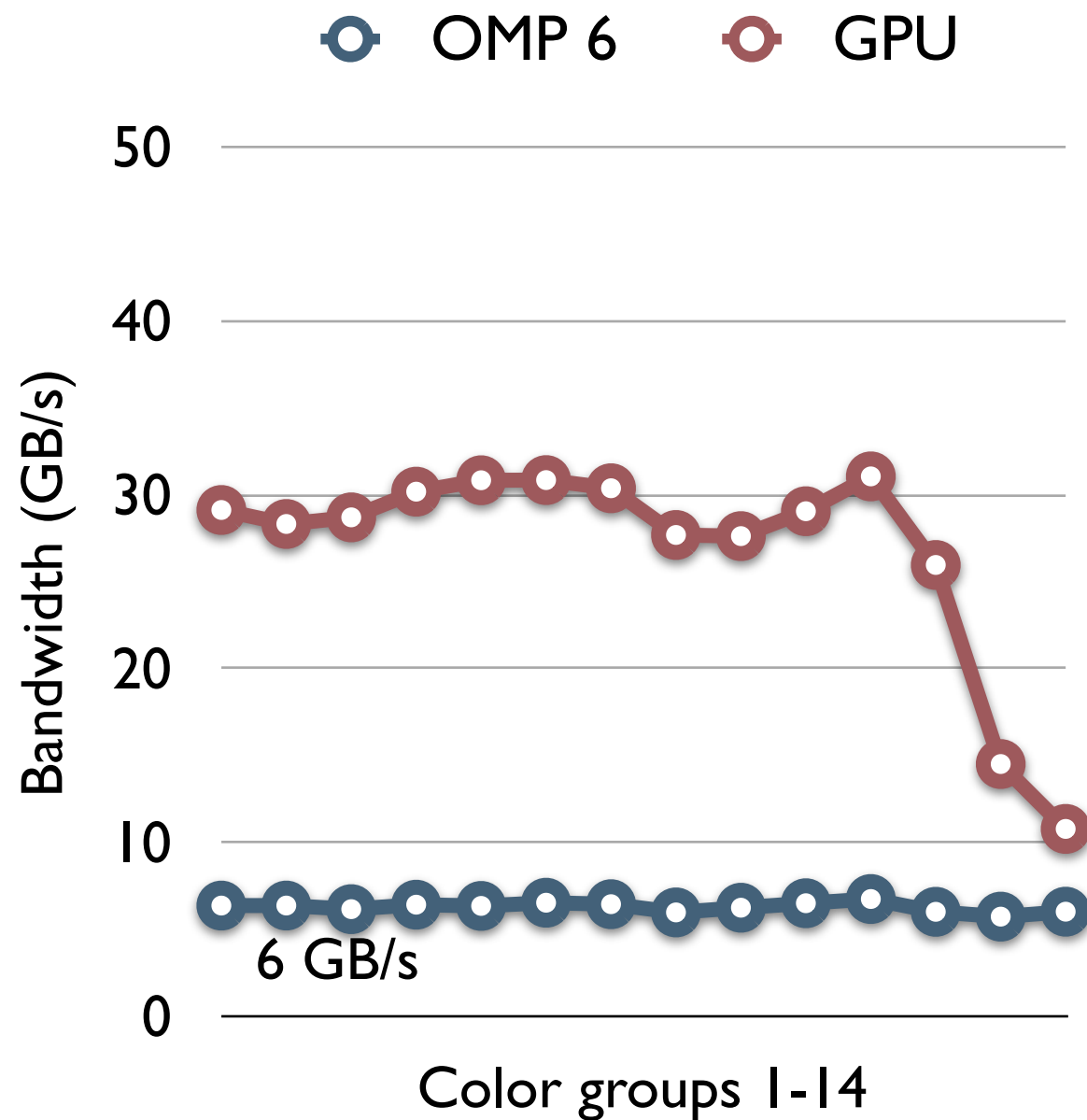
Kernel: inviscid fluxes with scalar dissipation in 2D



Efficiency of vector assembly (P2)



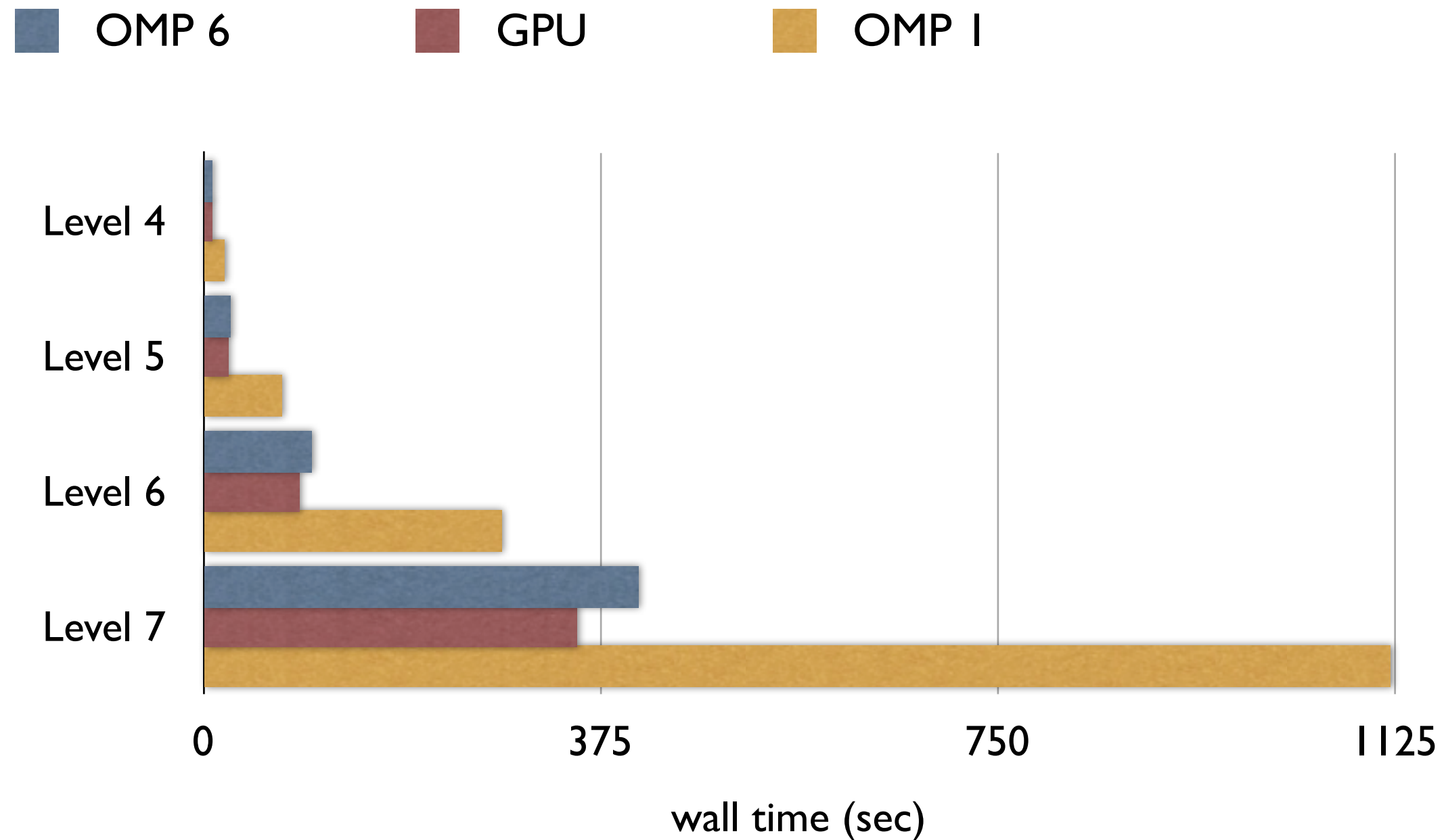
Kernel: inviscid fluxes with register intense Roe-type dissipation in 2D



Efficiency of „full“ simulation (P2)



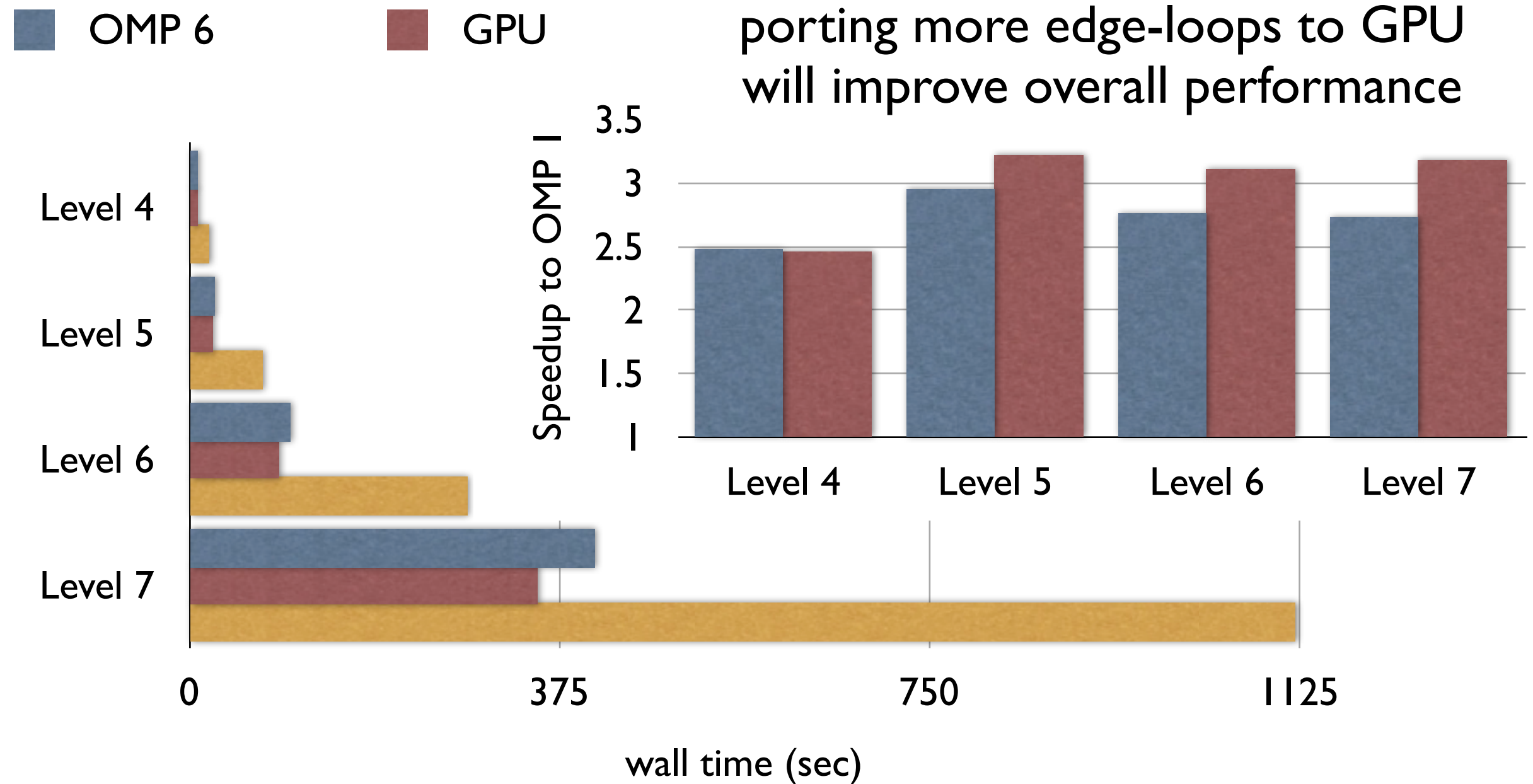
Comparison: full OpenMP vs. OpenMP + vector assembly on GPU



Efficiency of „full“ simulation (P2)



Comparison: full OpenMP vs. OpenMP + vector assembly on GPU

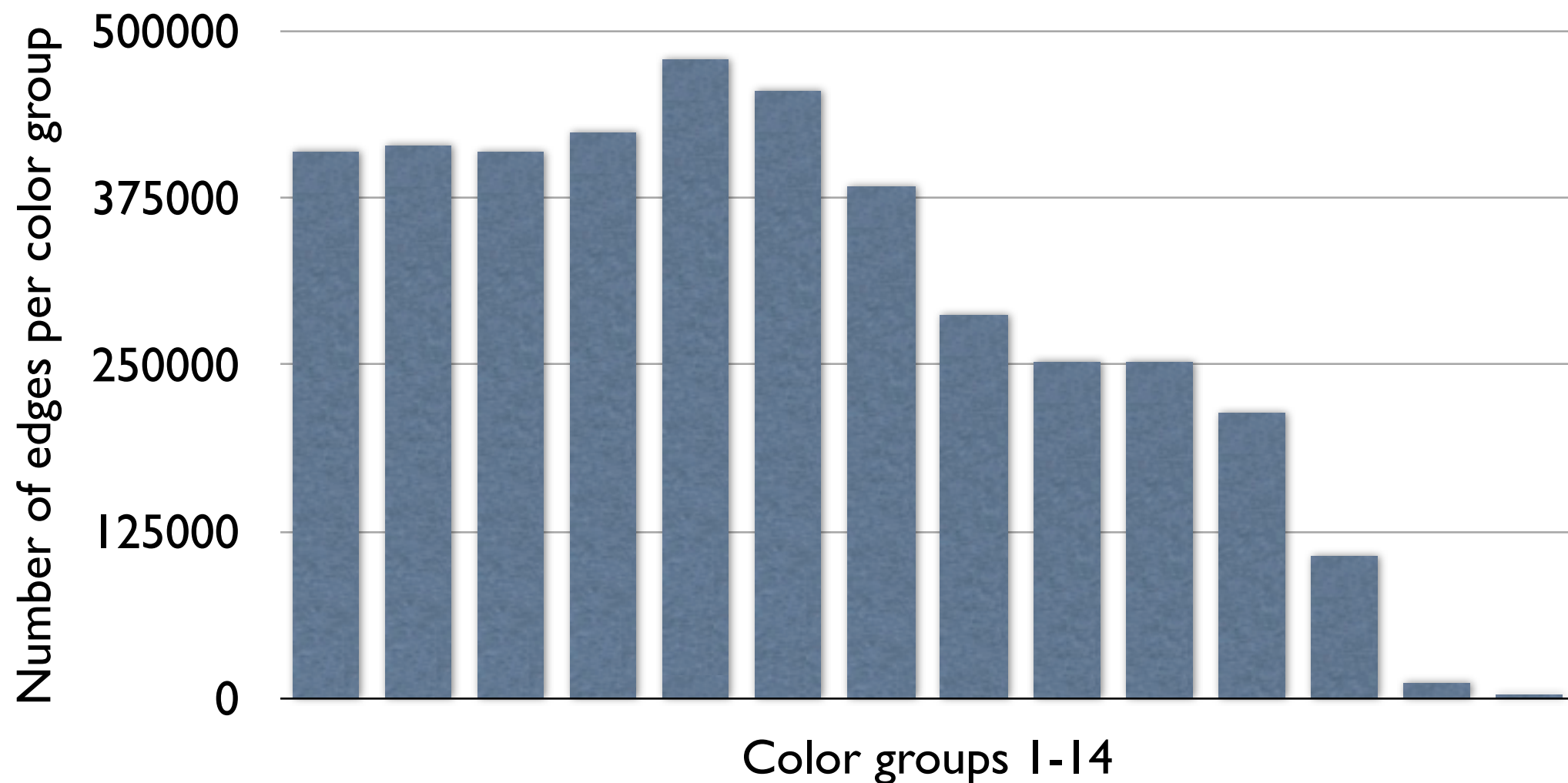


A second look on vector assembly



GPU would perform best for only few groups with
an equally distributed number of edges per color

=> use better coloring algorithm and „better“ finite elements (?)

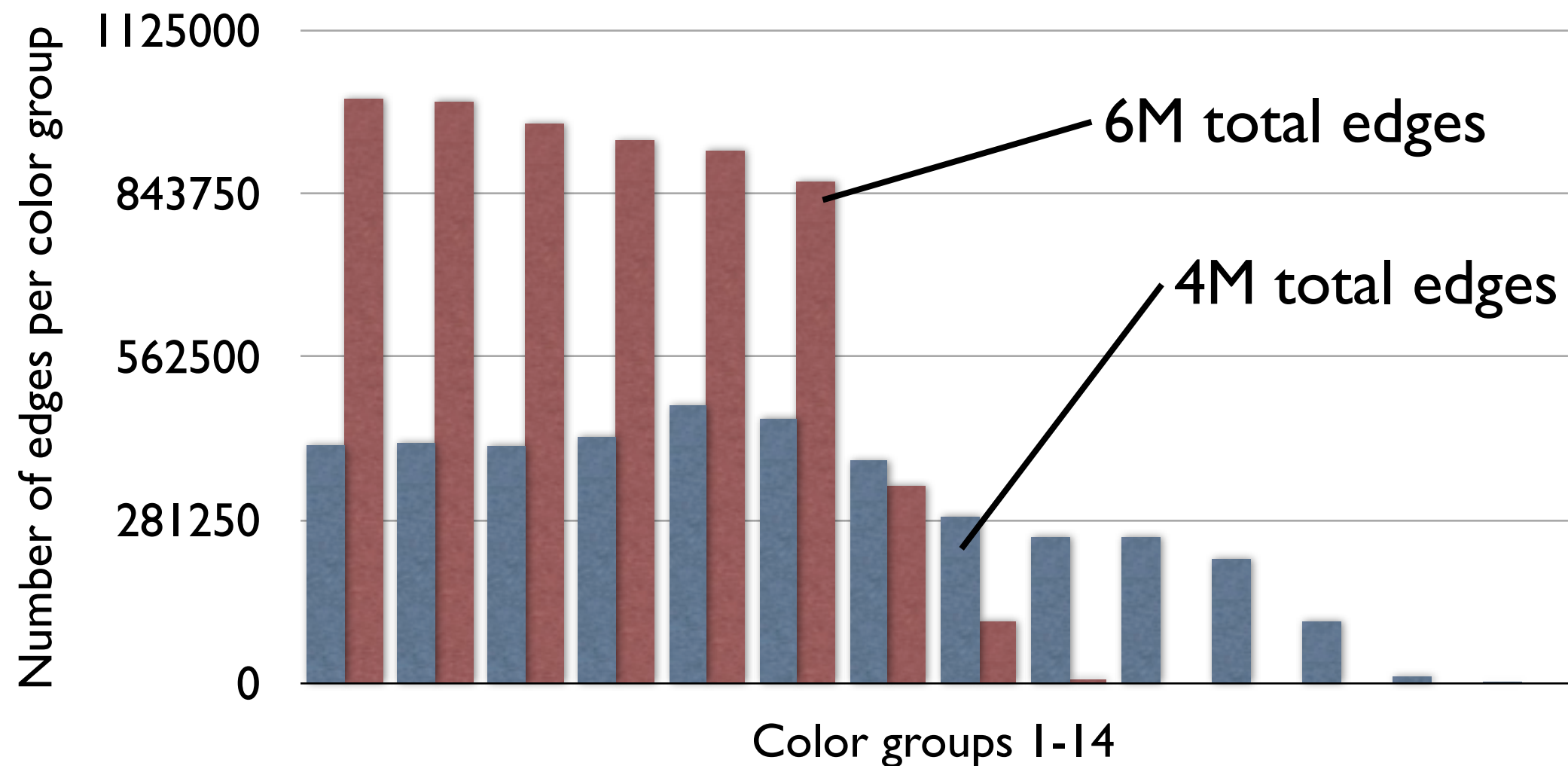
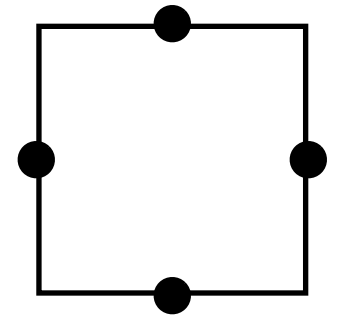


A second look on vector assembly



Non-conforming rotated bilinear Q1 ~ finite element

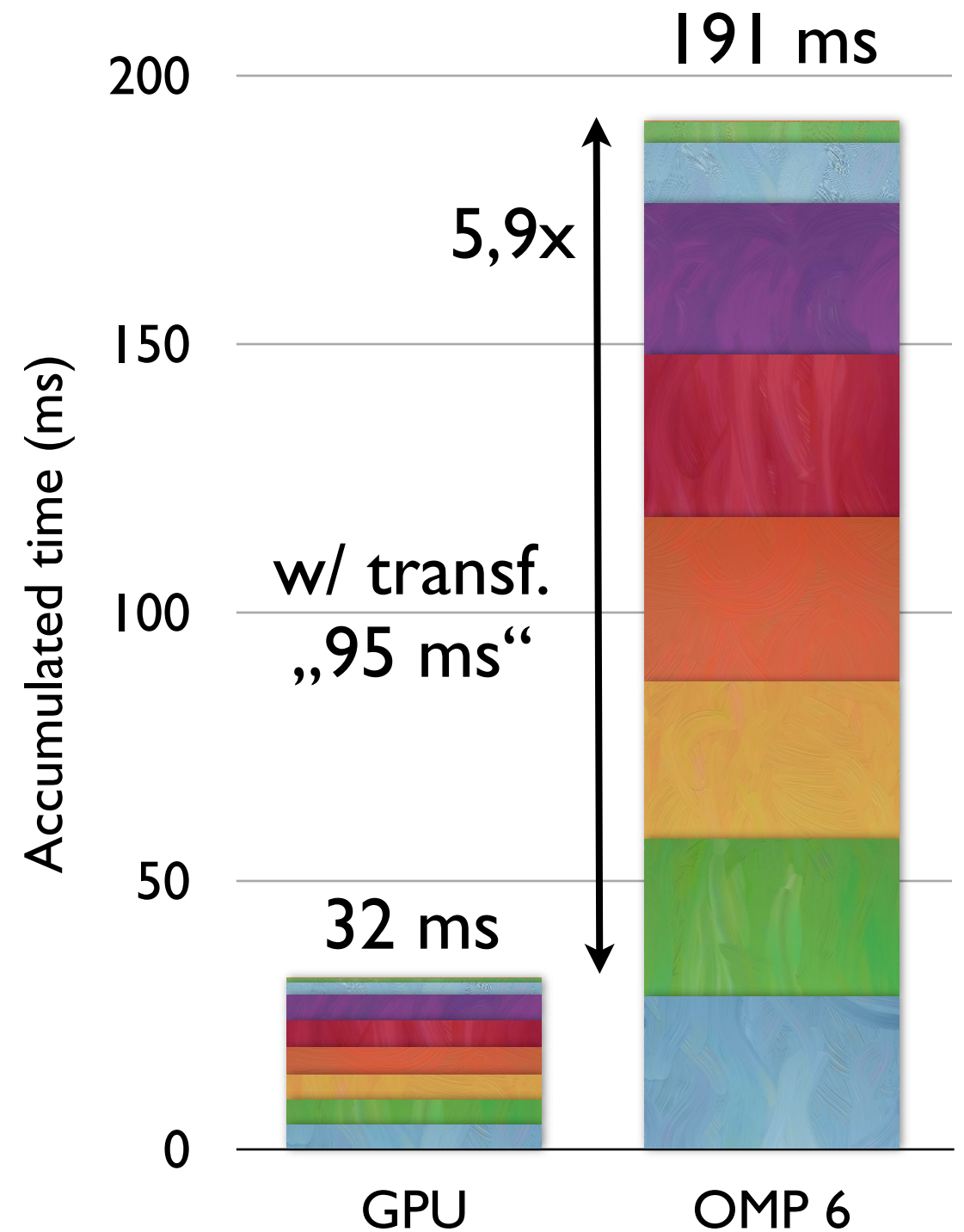
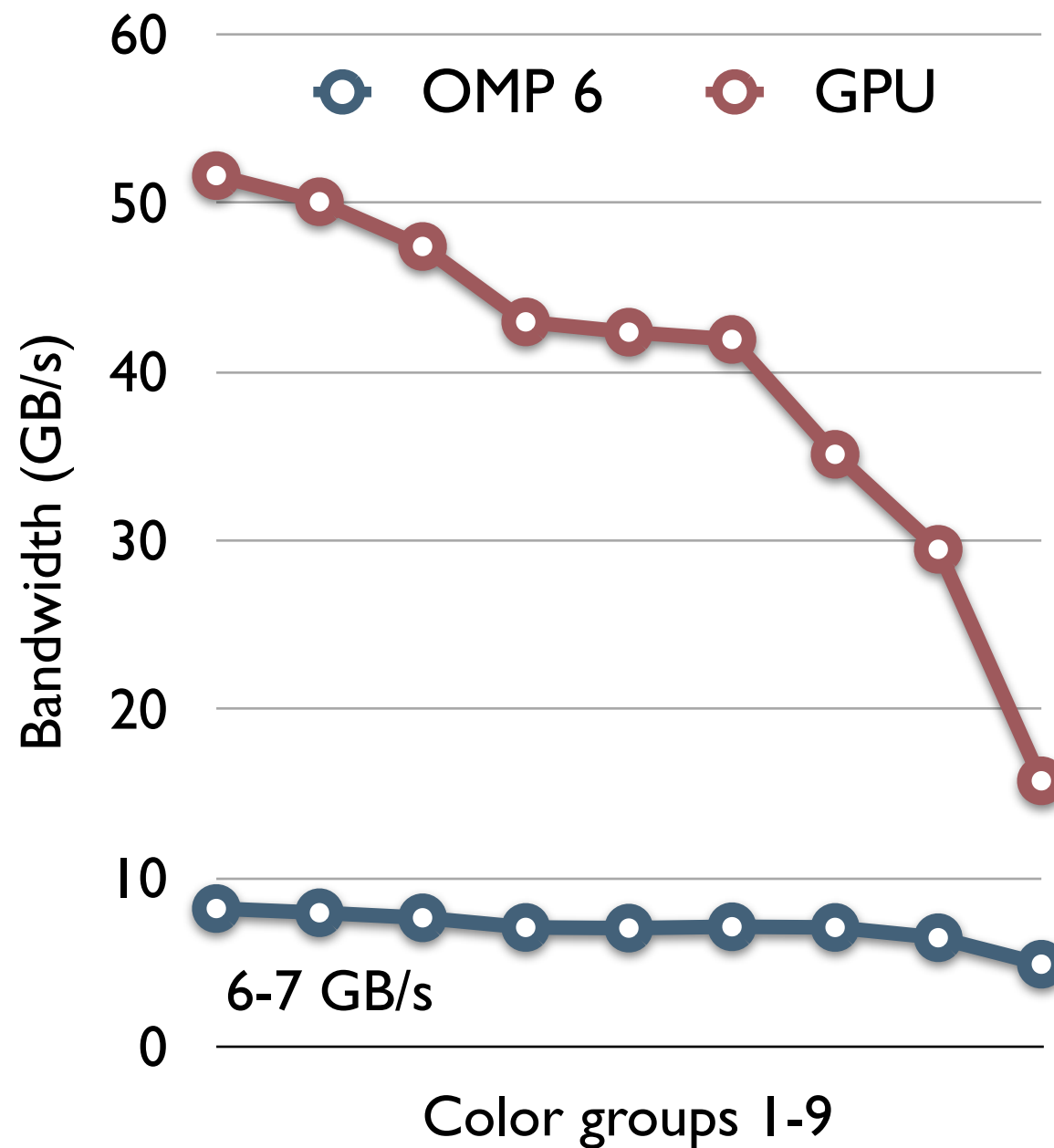
- DOFs are located at the midpoints of edges
- each DOF always couples with 6 neighbors (10 in 3D)



Non-conforming finite elements



Kernel: inviscid fluxes with scalar dissipation in 2D



Technical details

Handle-based storage manager

- (De-)allocate new memory assigned to ihandle (= unique integer)

call `storage_new((/a,b,c/), ST_DOUBLE, ihandle, [rheap])`

call `storage_free(ihandle, [rheap])`

- Associate 3D double pointer to memory at ihandle

call `storage_getbase(ihandle, p_Darray3D, [rheap])`

- Different memory managers can be used ,under the hood‘

Fortran only allocate / deallocate	CUDA <code>cudaHostAlloc</code> / <code>cudaFreeHost</code>
<ul style="list-style-type: none">■ works with all F90 compilers■ no asynchronous transfers	<ul style="list-style-type: none">■ works with F2003: <code>iso_c_binding</code>■ fast and asynchronous transfers

Handle-based storage manager, cont'd

- Full Fortran 95 functionality: `size`, `shape`, assumed-shape arrays

- Lightweight data structures (= collection of handles)

```
type t_matrix
```

```
integer :: na, neq, ncols
```

```
integer :: h_Da, h_Kcol, h_Kld    <- simple resize in mesh adaptation
```

- Co-Processor support does not break legacy code

- Memory transfers and accessing memory on device

```
storage_syncMem(ihandle, <direction>, async=YES/NO, ...)
```

```
storage_getMemPtr(ihandle, p_memptr, ...) -> pass p_memptr to GPU kernel
```

- Support for OpenCL, ... can be integrated easily (!?)

Meta-programming library

Example: Roe's Riemann solver is the same on CPU/GPU except for

- different programming languages (Fortran/C++)
- different index addressing (0-/1-based)
- different memory layouts (AoS/SoA/mixture)
- Meta-programming of application code via pre-processor macros
 - tedious to find least common subset of built-in pre-processor features supported by all Fortran compilers
 - GNU cpp: F90 \rightarrow f90 + Fortran compiler

Meta-programming library

```
IDX1( flux_xi, 1) = INVSCFLUX1_XDIR( edgedata, IDX3, i, tid, . . . )  
IDX1( flux_xi, 2) = INVSCFLUX2_XDIR( edgedata, IDX3, i, tid, . . . )  
...
```

#define FORTRAN_AOS

```
flux_xi(1) = edgedata(2, i, tid)  
flux_xi(2) = edgedata(2, i, tid)*ui+pi  
...
```

#define FORTRAN_SOA

```
flux_xi(1) = edgedata(tid, i, 2)  
flux_xi(2) = edgedata(tid, i, 2)*ui+pi  
...
```

#define C_AOS

```
flux_xi[((nthreads)*(1 +(-1))+(tid))] = edgedata[((nthreads)*(2)*(2 +(-1))  
                                                    +(nthreads)*(i +(-1))+(tid))];  
flux_xi[((nthreads)*(2 +(-1))+(tid))] = (edgedata[((nthreads)*(2)*(2 +(-1))+(  
                                                    (nthreads)*(i +(-1))+(tid))] *ui+pi);  
...
```

Summary

Parallelization of edge-based CFD-solver Featflow2

- Group-FEM formulation leads to memory and time efficient edge-based assembly on Many-Core architectures
- Minimally invasive integration of GPU acceleration in legacy code
- Meta-programming library simplifies mixing of programming languages and enables reuse of application code

Future plans

- Explore benefits of non-conforming FEs for edge-parallelization
- Port more edge-loops to CUDA and add multi-GPU support
- Combine assembly on CPU and GPU adaptively

References

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- D. Kuzmin, *Explicit and implicit FEM-FCT algorithms with flux linearization*. JCP 2009, 228(7), pp. 2517–2534.
- D. Kuzmin, M. M, S. Turek, *Multidimensional FEM-FCT schemes for arbitrary time-stepping*. IJNMF 2003, 42(3), pp. 265–295.
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- X. Roca, private communication (FEF 2011, Munich)